# DNA and Bernoulli Random Number Generator Based Security Key Generation Algorithm 

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#### Abstract

Security is a major concern for the communication sector. The technique presented in this paper provides a novel security key generation mechanism. The proposed technique aims to generate a security key using the biological characteristics of the human body and the mathematically generated pseudo random sequences, thus producing different keys for different individuals. The final key is produced through the fusion of deoxyribonucleic acid (DNA) sequence of 1024 characters and Bernoulli Random Number Generator sequence of 256 bits. The performance of produced keys is evaluated using National Institute of Standards and Technology (NIST) tests and uniqueness is verified through avalanche test.


Keywords: Authentication, Bernoulli random number generator, biometrics, communication, confidentiality, DNA, integrity, security

## INTRODUCTION

Security is critical in communication networks following the growth of the internet. Biometrics and integrity are two mechanisms which can improve security considerations. Biometrics ensures the identification and authentication of individuals by observing their personal unique features (Hao, Anderson, \& Daugman, 2006). A practical system that integrates the iris biometrics into cryptographic applications can be found in Hao et al. 2006. A system that works using audio fingerprints has also been proposed by (Brown \& Seberry, 2001; Chouakri, Bereksi, Ahmaidi, Fokapu, 2005; Covell \& Baluja, 2007. Ktata, Ouni, Ellouze, 2009) studied the use of electrocardiogram (ECG) signals for enhanced security. Baluja, Covell, 2007 proposed to create personal signatures for authentication. Chen \&

## Article history:

Received: 29 December 2016
Accepted: 21 April 2017

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Chandran, 2007) did a study on identification based on image processing is carried out by various researchers.

DNA has been used in many cryptographic algorithms to provide confidentiality (Chang, Kuo, Lo, \& Wei, 2012). Bernoulli random
number generator (BRNG) is used to make the technique more efficient and effective. The BRNG works on a secret seed value which ensures the generation of a different sequence for each seed value provided. The work reported in this paper is based on the idea of unique and random attributes of DNA and Bernoulli random number generator sequence.

The approach to generate a security key is implemented through numeric coding of the DNA sequence and the Bernoulli random sequence generated with the secret seed value. The output of key generating algorithm is tested using NIST tests of randomness as well as the strict avalanche criterion, the results of which are formulated in Table 4 and Table 5 respectively. Results point out the applicability of the proposed approach in areas where security is a concern.

The paper is organized as follows; characteristics of DNA and Bernoulli random number generator are described in Section 2. In Section 3, proposed algorithm for the 256bit key generation is presented where the DNA data is taken from the MIT-BIH database, followed by result and analysis in Section 4. Finally, conclusions and future work are reported in Section 5.

## Characteristics of DNA and BRNG

Progress in the field of biotechnology has made the DNA sequencing more effective. DNA sequencing for various organisms has been done with higher accuracy (Goldberger, Amaral, Glass, Hausdorff, Ivanov, Mark, Mietus, Moody, Peng \& Stanley, 2000). Investigating the biological relationships of different species is known as analysing the DNA sequence.

The internet provides various databases of DNA sequences which can be easily accessed from the World Wide Web (Ensembl Genome Browser, NCBI databases). The DNA is composed of two polymeric strands made of monomers that include a nitrogenous base (A-adenine, C-cytosine, G-guanine, and T-thymine), deoxyribose sugar and a phosphate group. According to most of the techniques a DNA sequence is taken as symbolic data that is composed of four characters $\mathrm{A}, \mathrm{C}, \mathrm{G}$ and T corresponding to the four types of nucleic acids present. The DNA samples from a genome are sequenced using a Genome Sequencer, the signals created in the sequencing process are then analysed by the software to generate millions of sequenced bases.

The backbone of each strand located on the surface of the DNA is formed by the sugar and phosphate groups, while the inside of the structure is made of bases. There exists a weak hydrogen bond between the complementary bases of each strand (i.e. between A and T and between C and G ) giving rise to pairing of bases, this pairing holds the two strands together. DNA sequences are unique for every individual, including the identical twins. The pattern formed by a DNA sequence specifically represents an individual and its characteristics. Hence, there is no possibility of duplicity in the DNA sequences.

To strengthen the bond of security, a random sequence is generated by Bernoulli random number generator. This sequence is generated with a seed value that is secretly given by the user. BRNG computes $\mathrm{n}^{\text {th }}$ Bernoulli number for a given integer n . Bernoulli numbers are a sequence $B_{n}$ of rational numbers defined by the Taylor expansion.

Bernoulli numbers have a prominent place in mathematics, for instance they appear in Taylor expansion of tangent and in Euler-Maclaurin formula. Equation (1) represents the BRNG method for generation of random sequences.

$$
\begin{equation*}
\sum_{n=0}^{\infty} \frac{B_{n}}{n!} x^{n}=\frac{x}{e^{x}-1} \tag{1}
\end{equation*}
$$

The fusion of BRNG and DNA sequence forms a very strong 256-bit key which is less susceptible to attacks and thus provides higher level of security.

## The Key Generation Mechanism

The proposed key is generated by combining DNA sequence of an individual and the random sequence obtained through BRNG using a secret seed value. The method of the generation of these keys is devised into three subsequent subsections and the hierarchical structure is portrayed in Figure 1.

## DNA sequence formulation

1) Obtaining a DNA sequence of 1024 characters from the DNA database [9]: The DNA sequence consists of 1024 characters forming base pairs 'agct'.
2) Obtaining the binary sequence from DNA characters:

Each character of the DNA sequence is represented by 8 bit ASCII code. Hence, resulting in a DNA sequence of length 8192 bits.
3) Framing a DNA sequence of 256 bits:
(i) The DNA sequence is then divided into equal halves.
(ii) Apply exclusive-or operation on the obtained sequences.
(iii) The result is further divided into two equal parts and exclusive-or operation is applied again.
The step (iii) is repeated till a sequence is obtained whose length is 256 bits. The whole procedure is summarized in the flow chart (Figure 1).

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Figure 1. Key generation process

The conversion from 'agct' sequence of 1024 characters to 256 -bit binary DNA sequence is given in Table 1.

Table 1
DNA sequence formulation

|  | DNA sequence (1024 characters) | DNA sequence (8192 bits) | DNA sequence (256 bits) |
| :---: | :---: | :---: | :---: |
| $\mathrm{D}_{1}$ | gcacaatcagaagcaggcggaggagacggcg gccttcgaggaggtcatgaaggacctgagcctg ac. $\qquad$ $\qquad$ $\qquad$ $\qquad$ $\qquad$ gcaca <br> gaggcaaggcgtcagcaggcatcgcccaccet gtctccgctgtcacccatcactcaggctgtagcc atg. |  | ```000001000000000000001010100 01011100000010000001100001 01010000010000000110000001 10000001100000011000010011 000000000000000000001001100 00000000000010000001100000 00000000011000000110000001 00000000100001001100000110 000000100000000000000010000 0000000001001100000000``` |
| $\mathrm{D}_{2}$ | ccacgegtccgggegagaagatggcgacttcg aacaatccgcggaaattcagcgagaagatcgc gct. $\qquad$ $\qquad$ $\qquad$ $\qquad$ $\qquad$ ccctgtccetgagcacagaggcaaggcgtcag caggcatcgccegccetgtccccgctgtcacce at. |  | 00000010000000100000001000 00001000000110000101110000 01000001011100000100000000 00000000000000000000000110 00000010000001100001010100 01001100000010000000000000 00000000001000010011000001 10000000000000001000000110 00000110000000100000000000 0101110000010000000100 |
| $\mathrm{D}_{3}$ | agccettaggggaagagtcctgctctggctgttg atgctccagctccagaaatcccagtacctgcaac tg. $\qquad$ $\qquad$ $\qquad$ $\qquad$ ccctacgecttcttcacceaggegggctcccagc agccaccgccgcagccccagcccccgccg |  | 00010101000000000000001000 01011100000100000101110000 00000000001000000100000001 10000000000000010000010101 00000010000001000000000000 00000000010101000000100000 00100000000000000100000000 00000101110000001000010011 00010011000001100000000000 0101110001001100010011 |

[^0]Random sequence generation through BRNG. Three sequences are generated using three different seed values. A seed value is a secret input given by the user at the beginning of generator operation. The obtained random sequences along with the corresponding seed value are summarized as in Table 2.

Table 2
BRNG sequence
Seed value Bernoulli random sequence (256 bits)
$B_{1} 1231 \quad 1011001100000011010110100110000101010010000101110111101110011101000$ 0000110000011010101101000101011011101010001101110101010101100000001 1001110110100000110110011001011110111001010001000101100100000000101 1100001111011010111011110000011011110000110010010010110
$B_{2} 10135500010110100100010101000110010110011100000011010111111001111010101110$ 00100101100101000011101101110011100001001100110011011011101001011100 110111010001100000110111110011111111100110100000111110011001001001100 010001000110011011100100100100101010110101101010001
$B_{3} 2114 \quad 0111010100001101001100000101010100000001000100001010101010101000001$ 10101111010101010011010011010010100000110001000101011100000001000111 00111000100001100100101111100010111000101001101011101111110101100011 00011000101110110010110100100100001111011110001101011
${ }^{*} \mathbf{B}_{1}, \mathbf{B}_{2}, \mathbf{B}_{3}$ : Three BRNG sequences

Fusion of DNA sequence and BRNG sequence. To enhance the strength of security and randomness in the key, exclusive-or logic is applied between each DNA and BRNG sequence. The resultant key has length of 256 bits and is strong enough to survive critical attacks on networks. The process is repeated for two other DNA and BRNG sequences. Therefore, three 256 bit keys are obtained as shown in Table 3.

Table 3
Security keys

KEY $_{1} \quad 00010010100100010100010010000001011100100011001111101100111011101110010001$ 01111101000101101100010010101101001100110011011010100101011100110111110001 111000110111110010011111111110100100111110111000000101100100001000010011011 100100000100101010111100001010001

## KEY $_{2} 10110001000000010101100001100010101011110001110010111001110110011000010011$ 0000011010101101000101011010001010000101110011010000110001000000111001010 00001101100110010110101100001100011101011001000000011011101101111000010111 00111000001101010110011011001001111

$\mathrm{KEY}_{3} \quad 0110000000001101001100100100001000000101000001111010101010101010001100011$ 11011001010011010011110010001010110000000101111100000001000111001100100000 01110100101011100010111000001001101011100100010101110011100000000010001100 01110100100100100001010111101111000

The final key is a combination of true random source (DNA) and pseudo random source (BRNG). The keys obtained have unique and random characteristics and thus can be useful in maintaining security.

## RESULTS AND DISCUSSION

A security key is said to be efficient if it is random and unique. The National Institute of Standards and Technology (NIST) tests discuss some aspects of selecting and testing random number generators. The outputs of these generators may be used in most of the security applications for the generation of security keys. The generators that are to be used for security applications should meet stronger requirements than those to be used for other applications. To be precise, their outputs need to be unpredictable for unknown inputs. These tests may be useful for determining if a generator is suitable for a particular security application. The randomness of a key is evaluated on the bases of its P -value, which must be greater than 0.01 for a random sequence.

## NIST Tests

The keys obtained are evaluated on the basis of seven NIST randomness verification tests and the P-value is calculated for each test with respect to the security key. An overview of the NIST tests used for evaluating the sequences is given as:

Runs test. The purpose of applying this test is to calculate the number of runs in an entire sequence, where run specifies the number of uninterrupted sequence of identical bits (Rukhin et al., 2010). The results in Table 4 clearly depict that the proposed algorithm has higher rate of interruptions.

Frequency Test. Frequency test is used for determining the proportion of number of ones and zeros in an entire sequence. It is used to check the closeness between the number of ones and number of zeros. A sequence is said to be random if the proportion of zeros and ones are close to each other (Rukhin et al., 2010). The results in Table 4 clearly depict that the proposed algorithm produces better proximity between the count of ones and zeros as compared to the traditional techniques.

Approximate Entropy Test. This test deals with the frequency of all the overlapping bit patterns across the entire sequence. The aim of this test is to compare the frequency of overlapping blocks of two consecutive/adjacent lengths with the expected result for a random sequence.

Discrete Fourier Transform Test (DFT). The DFT test is used to find the peak heights in the Discrete Fourier Transform of a sequence. The aim of this test is to detect periodic features (i.e. repetitive patterns) in the tested sequence that indicates a deviation from the assumed randomness. The goal is to detect if the number of peaks exceeding the $95 \%$ threshold are significantly different than $5 \%$.

Binary Derivative Test. The test is performed using exclusive-or operation between successive bits until only one bit is left. Afterwards, the ratio of number of ones to the length of entire sequence in each case is calculated. Finally, the average of the ratio of all the sequences is observed, and where the value lies near to 0.5 , then the sequence is considered as a random sequence (Rukhin et al., 2010). The results in Table 4 show that the output of proposed algorithm is random.

Maurer's "Universal Statistical" Test. The aim of this test is to detect if a sequence can be significantly compressed without any loss of information. The number of bits between matching patterns are calculated. A sequence that is significantly compressible is considered to be nonrandom. This test is also known as Universal test (Rukhin et al., 2010).

Random Excursion Variant Test. The test focuses on the total number of times a particular state occurs in a cumulative sum random walk. It detects the deviations from the expected number of visits. The P -value specifies if the sequence is random or not. This test considers successive sums of the binary bits as a one-dimensional random walk (Rukhin et al., 2010).

$$
\begin{equation*}
P_{\text {value }}=\operatorname{erfc} \times \frac{(|\xi(x)-j|)}{\sqrt{(2 \times j \times((4 \times|x|)-2))}} \tag{1}
\end{equation*}
$$

Where,
erfc: the error function
$\xi$ : the total number of times the state $x$ occurs
$x$ : the state occurred
$j$ : the total number of cycles
The P-value must be greater than 0.01 for a sequence to be random (Rukhin et al., 2010).
It is observed that the proposed technique produces better results in terms of randomness of the keys. The efficiency of the proposed technique is evaluated by comparing it with other traditional biometric techniques used for authentication and security key generation. The tests have been performed on $\mathrm{KEY}_{1}$ and results are tabulated in Table 4.

Table 4
Comparison and analysis

| Input Source of random number generator | Key <br> Length (bits) | Runs <br> Test | Frequency Test | Approximate Entropy Test | DFT Test | Binary Derivative Test | Maurer's <br> Test | Random <br> Excursion <br> Variant |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | P-value | P-value | P-value | P-value | P-value | P-value |  |
| 1. ECG | 128 | 0.1262 | 0.2487 | 0.5468 | 0.0294 | 0.5039 | 0.9428 | Random |
| 2. Image | 256 | 0.0809 | 0.8026 | 0.9759 | 0.4220 | 0.4887 | 0.9780 | Random |
| 3. Iris sequence | 128 | 0.1254 | 0.3768 | 0.9409 | 0.3304 | 0.5021 | 0.9062 | Random |
| 4. Finger print | 128 | 0.3345 | 0.3041 | 0.3345 | 0.7597 | - | 0.2757 | Random |
| 5. DNA \& BRNG | 256 | 0.0438 | 0.9005 | 0.9340 | 0.8185 | 0.5090 | 0.9920 | Random |

It can be observed that the P-value obtained for the keys generated for all the seven tests is significantly greater than 0.01 . Thus, it can be concluded that the keys generated are random in nature and hence fulfil the basic criteria required for security keys.

Avalanche test has also been performed on the keys obtained. The purpose of this test is to check the avalanche effect, which is a desirable property of the security keys. Wherein if the input is changed slightly the output changes significantly. It gives the percentage of bits flipped with a change in the input. It is a desirable property of security keys.

The test is performed on three sets of DNA and BRNG sequences:
Case 1: In the initial set, two security keys are generated through two DNA sequences while keeping the same BRNG sequence.
Case 2: The second set involves generation of two security keys through the same DNA sequence and two BRNG sequences.
Case 3: In the third set, two security keys are generated through two DNA and BRNG sequences.
Further, the avalanche effect is calculated for each of the three sets and results are tabulated in Table 5, Table 6 and Table 7 respectively, this is done to know the amount of randomness the proposed technique produces on changing the input.

The avalanche effect can be calculated using the formula given in equation (3).

$$
\begin{equation*}
\text { Avalanche effect }=\frac{\text { No.of bits flipped in the sequence }}{\text { Total no.of bits in the sequence }} \times 100 \tag{3}
\end{equation*}
$$

The result of avalanche effect on changing the inputs is summarized in tabular form, where $D_{1}, D_{2}, D_{3}$ represent the DNA sequences as taken from Table 1 and $B_{1}, B_{2} B_{3}$ represent the BRNG sequences characterized by the seed values, as taken from Table 2.

Table 5
Avalanche test analysis: Case 1

| DNA <br> Sequences $\left(D_{n}\right)$ | Seed <br> Value | Bernoulli <br> Random <br> Sequence $\left(\mathbf{B}_{\mathrm{n}}\right)$ | Key Generated$K=D_{n} \operatorname{xor} B_{n}$ | Avalanche Result of Key (K) |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | No. of bits flipped | Avalanche Effect |
| $\mathrm{D}_{1}$ | 1231 |  | 1011011100000011010011110111011001 |  |  |
|  |  |  | 010000000100010110111010011001000 |  |  |
|  |  |  | 001111000010101010000100011001100 |  |  |
|  |  |  | 1110010001101110101010111111000001 |  |  |
|  |  |  | 100111010010000101011001100101100 |  |  |
|  |  |  | 011100011000101010110011000010001 |  |  |
|  |  |  | 1110011111101111011101111000011101 |  |  |
|  |  |  | $1110000111011110010110$ |  |  |
| $\mathrm{D}_{2}$ |  | $\mathrm{B}_{1}$ | 101100010000000101011000011000110 | 58 | 22.65 \% |
|  |  |  | 101010000000000011111111000101000 |  |  |
|  |  |  | 000101100000110101011010001010110 |  |  |
|  |  |  | 110110100010011101100101110010001 |  |  |
|  |  |  | 010101110100100000110110011001011 |  |  |
|  |  |  | 100111101100001011101100100000000 |  |  |
|  |  |  | 0011100111111010110111010110000011 |  |  |
|  |  |  | 011011110110000010010010 |  |  |

[^1]Table 6
Avalanche test analysis: Case 2

| DNA <br> Sequence | Seed <br> Value | Bernoulli Random | Key Generated | Avalanch Key | Result of (K) |
| :---: | :---: | :---: | :---: | :---: | :---: |
| ( $\mathrm{D}_{\mathrm{n}}$ ) |  | $\begin{aligned} & \text { Sequence } \\ & \left(\mathbf{B}_{\mathrm{n}}\right) \end{aligned}$ | $K=D_{n}$ xor $B_{n}$ | No. of bits flipped | Avalanche effect |
|  |  |  | 101101110000001101001111011101100 |  |  |
|  |  |  | 101000000010001011011101001100100 |  |  |
|  |  |  | 000111100001010101000010001100110 |  |  |
|  |  |  | 011100100011011101010101111110000 |  |  |
|  | 1231 | $\mathrm{B}_{1}$ | 011001110100100001010110011001011 |  |  |
|  |  |  | 000111000110001010101100110000100 |  |  |
|  |  |  | 0111100111111011110111011110000111 |  |  |
|  |  |  | 011110000111011110010110 | 124 | 48.43 \% |
| $\mathrm{D}_{1}$ |  |  | 000100101001000101000100100000010 |  |  |
|  |  |  | 1110010001100111110110011101110111 |  |  |
|  |  |  | 0010001011111010001011011000100101 |  |  |
|  | 101355 | $\mathrm{B}_{2}$ | 0110100110011001101101010010101110 |  |  |
|  |  |  | 0110111110001111000110111110010011 |  |  |
|  |  |  | 1111111101001001111101110000001011 |  |  |
|  |  |  | 0010000100001001101110010000010010 |  |  |
|  |  |  | 1010111100001010001 |  |  |

${ }^{*}$ Refer Table 1 for $\mathrm{D}_{1,} \mathrm{D}_{2}, \mathrm{D}_{3}$ and Table 2 for $\mathrm{B}_{1}, \mathrm{~B}_{2}, \mathrm{~B}_{3}$
${ }^{* *}$ Same DNA sequences - Different BRNG sequences

Table 7
Avalanche test analysis: Case 3

| DNA <br> Sequence $\left(D_{n}\right)$ | Seed <br> Value | Bernoulli Random Sequence ( $\mathrm{B}_{\mathrm{n}}$ ) | Key Generated$K=D_{n} \operatorname{xor} B_{n}$ | Avalanche Result of Key <br> (K) |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | No. of bits flipped | Avalanche Effect |
| $\mathrm{D}_{1}$ | 1231 | $\mathrm{B}_{1}$ | 10110111000000110100111101110110 | 58 | 22.65 \% |
|  |  |  | 01010000000100010110111010011001 |  |  |
|  |  |  | 00000111100001010101000010001100 |  |  |
|  |  |  | 11001110010001101110101010111111 |  |  |
|  |  |  | 00000110011101001000010101100110 |  |  |
|  |  |  | 01011000111000110001010101100110 |  |  |
|  |  |  | 000100011110011111101111011101111 |  |  |
|  |  |  | 000011101111000011101111001011011 |  |  |
| $\mathrm{D}_{2}$ | 101355 | $\mathrm{B}_{2}$ | 00010100100100110101001110010100 |  |  |
|  |  |  | 011101100010001011111101111111011 |  |  |
|  |  |  | 110011001011001010000111011011100 |  |  |
|  |  |  | 111110010011101100101110101111010 |  |  |
|  |  |  | 011111101111100011000001101111100 |  |  |
|  |  |  | 110111101010101001101111100110010 |  |  |
|  |  |  | 00001100100001001010011010100100 |  |  |
|  |  |  | 100100000100110111101010101 |  |  |

[^2]The avalanche test results indicate that a slight change in the inputs leads to a significant change in the output. The higher the percentage value of avalanche effect, the more is the uniqueness of the key. As the seed value may vary according to the user's choice and the DNA being unique to an individual, it can be assumed that the security system built has higher efficiency and is less susceptible to attacks.

## CONCLUSION

In this work, a novel algorithm is proposed to generate biometric security keys for authentication using an individual's DNA sequence and BRNG sequence. DNA being the most unique characteristic of an individual when collaborated with the Bernoulli seed is more efficient. The proposed key finds its application in various areas where maintaining the sender's confidentiality as well as retaining the original message are important. Most of the security system researches have been carried out on fingerprint, facial, iris and voice recognition; while those focused on the DNA are rare. A system that computes a 256 -bit biometric security key through the fusion of DNA and BRNG is presented and analysed using NIST Tests. Results indicate this method is superior to more traditional techniques.

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[^0]:    ${ }^{*} \mathbf{D}_{1}, \mathbf{D}_{\mathbf{2}}, \mathbf{D}_{3}$ : DNA sequences, NCBI Database

[^1]:    ${ }^{*}$ Refer Table 1 for $D_{1}, D_{2}, D_{3}$ and Table 2 for $B_{1,}, B_{2}, B_{3}$
    ${ }^{* *}$ Different DNA sequences - Same BRNG sequence

[^2]:    *Refer Table 1 for $D_{1,} D_{2}, D_{3}$ and Table 2 for $B_{1,} B_{2}, B_{3}$
    ${ }^{* *}$ Different DNA sequences - Different BRNG sequences

